

Binary Search: Search method that works on sorted sets of objects, typically stored within an array, to find the location matching the search key

Basic procedure checks the middle value:

If value matches, return location.

If key is greater than value, we can reduce search space to the upper half

If key is less than value, we can reduce search space to the lower half

Example:

[10   11   12   100   200   201   304   500   501   499]	key = 12
0    1    2    3    4↑    5    6    7    8    9 middle	

$12 < 200 \Rightarrow$  value must be located in lower half ( $0 \rightarrow 3$ )

\* Process repeats until key is found, or there are no more elements to search

\* Need a method to keep track of the elements we are searching.

↳ can use a left and right index to keep track of current search space

① Init left  $\rightarrow 0$  and right  $\rightarrow \text{size} - 1$

② while there are elements to search  $\equiv$  while  $\text{left} \leq \text{right}$

a)  $\text{middle} \rightarrow (\text{left} + \text{right}) / 2$

b) if  $\text{data-vals}[\text{middle}] == \text{key}$ , return  $\text{middle}$

c) if  $\text{data-vals}[\text{middle}] < \text{key}$ ,  $\text{left} \rightarrow \text{middle} + 1$

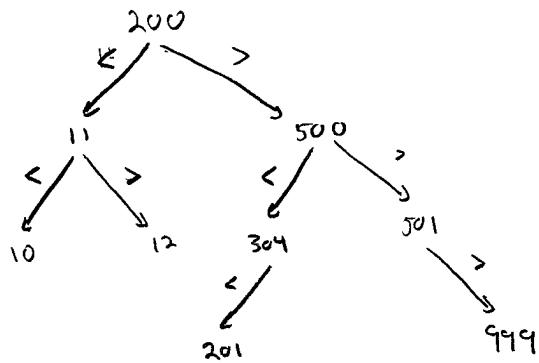
d) else  $\text{right} = \text{middle} - 1$

③ return -1

\* What is the complexity?

## Search Pattern Example:

[ 10 | 11 | 12 | 200 | 201 | 304 | 500 | 501 | 999 ]



\*What if we are using some data that is not stored in an array?

↳ Can we create a data structure to store elements/nodes and support the addition and deletion of nodes?

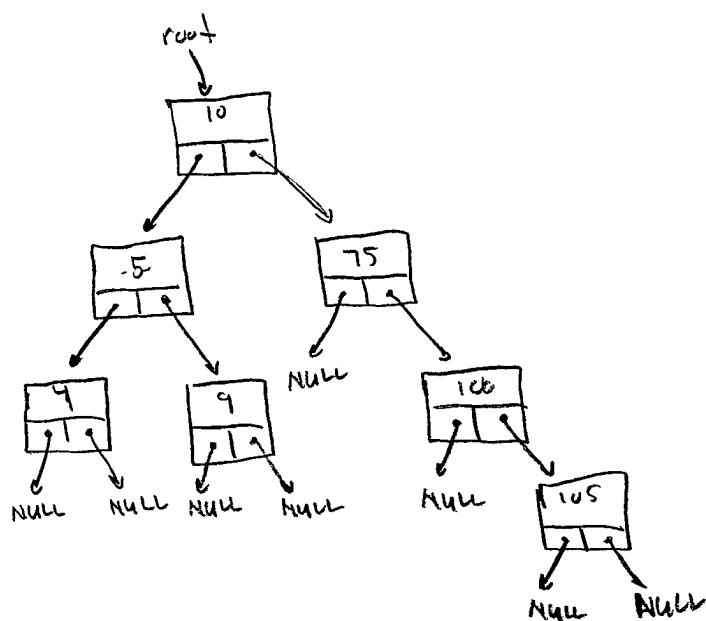
## Binary Trees:

- Hierarchical arrangements of nodes in which each node can have two nodes immediately below it
- Each node consists of data and two pointers to nodes one level below
- Nodes one level below current node are called children / descendants
- Node one level above current node is called parent
- Left and right pointers typically used to represent pointers within node



- Node with no children is a leaf node
- Root node is the single node at the top level of hierarchy

### Example:



Note: Nodes could also contain pointer to parent node.

## Binary Tree Declarations:

```
typedef struct BiTreeNode {
    int data;
    struct BiTreeNode * left;
    struct BiTreeNode * right;
} BiTreeNode;
```

```
typedef struct BiTree {
    BiTreeNode * root;
    int size;
} BiTree;
```

```
void bitree_init (BiTree *tree);
void bitree_destroy (BiTree *tree);
int bitree_ins_left (BiTree *tree, BiTreeNode *node, int data);
int bitree_ins_right (BiTree *tree, BiTreeNode *node, int data);
int bitree_rem_left (BiTree *tree, BiTreeNode *node);
int bitree_rem_right (BiTree *tree, BiTreeNode *node);
```

Binary Search Tree: Binary tree in which nodes are organized to aid in efficient searching

- An element within the node is used as a key to determine how nodes are organized
- All elements within the left subtree will have a smaller key than the current node
- All elements within the right subtree will have a larger key than the current node
- Duplicate keys are not allowed.

What is the complexity of a binary search tree search?

Binary Search Tree Interface:

```
int bstree_insert (Bstree *tree, int val);
int bstree_remove (Bstree *tree, int val);
BstreeNode *bstree_lookup (Bstree *tree, int search-key);
```

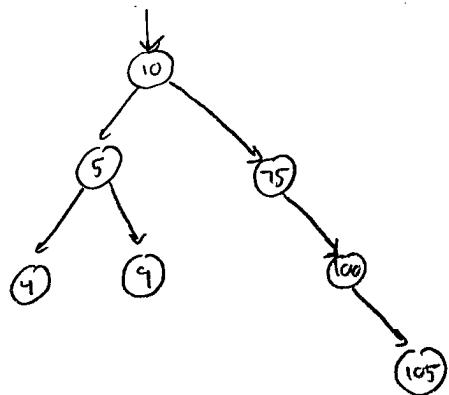
## Tree Traversals:

Preorder: Traverse root, Traverse left, Traverse right

Postorder: Traverse left, traverse right, traverse root

Inorder: Traverse left, traverse root, traverse right

### Example: Printing Tree



Preorder: 10 5 4 9 75 100 105

Postorder: 4 9 5 105 100 75 10

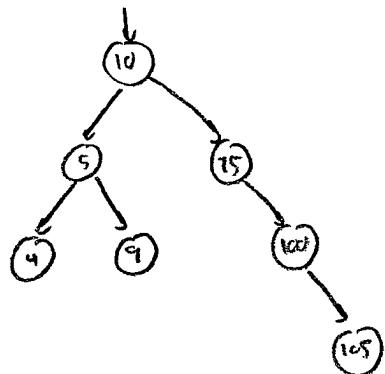
Inorder: 4 5 9 10 75 100 105

## Inorder Binary Tree Traversal (Recursion)

```
void bitree-print-inorder (BiTreeNode *node);
```

```
if node != NULL then
    bitree-print-inorder (node->left).
    print node->val
    bitree-print-inorder (node->right)
endif
```

Example:



Note: Integer values below used to illustrate current node

## bitree-print-inorder (10)

↳ if 10 != NULL then

i      bitree\_print\_order(10 → left ≡ 5)

↳ f 5 != NULL then

bitree-print-order ( $s \rightarrow \text{left} \equiv 4$ )

↳ if  $y \neq \text{NULL}$  then

bitree\\_print\\_morder ( $4 \rightarrow \text{left} \equiv \text{Null}$ )

↳ if null != null then

encl: 4

←  
print 4

bitree-print-Morder ( $\forall \rightarrow \text{right} = \text{NULL}$ )

↳ if null := null then

1

end

← print 5

bitree-print-order ( $5 \rightarrow \text{right} = 9$ )

↳ if  $q \neq \text{NULL}$  then

bitree-print-order( $q \rightarrow \text{left} \equiv \text{null}$ )

↳ if null != null then

卷之三

9

bitree print - inorder ( $q \rightarrow right \equiv \text{NULL}$ )

If null != null then

卷之三

1

end

print 10

bitree-print-node (10 → right ≡ 75)

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