ECE 274 Digital Logic - Fall 2008

RTL Design - Determining Clock Frequency and Behavioral RTL Design: C-to-Gates

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## Digital Design

Chapter 5:

## RTL Design

Slides to accompany the textbook Digital Design, First Edition, by Frank Vahid, John Wiley and Sons Publishers, 2007. htpp://www.dvvahid.com


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## RTL Design

Determining Clock Frequency

- Designers of digital circuits often want fastest performance
- Means want high clock frequency
- Frequency limited by longest register-to-register delay
- Known as critical path
- If clock is any faster, incorrect data may be stored into register
$\square$ Longest path on right is 2 ns

- Ignoring wire delays, and
register setup and hold times, for simplicity


## RTL Design <br> Critical Path

- Example shows four paths
- a to c through +: 2 ns
$\square$ a to d through + and *: 7 ns
- b to d through + and *: 7 ns
- b to d through *: 5 ns
- Longest path is thus 7 ns


## - Fastest frequency

口 $1 / 7 \mathrm{~ns}=142 \mathrm{MHz}$


RTL Design
Critical Path Considering Wire Delays

- Real wires have delay too
- Must include in critical path
- Example shows two paths
- Each is $0.5+2+0.5=3 \mathrm{~ns}$
- Trend
- 1980s/1990s: Wire delays were tiny compared to logic delays
- But wire delays not shrinking as fast as logic delays
- Wire delays may even be greater than logic delays!
- Must also consider register setup and hold times, also add to path
Then add some time to the computed path, just to be safe
- e.g., if path is 3 ns , say 4 ns instead


## RTL Design <br> A Circuit May Have Numerous Paths

## - Paths can exist

- In the datapath
- In the controller
- Between the controller and datapath
- May be hundred or thousands of paths
- Timing analysis tools that evaluate all possible paths automatically very helpful



## RTL Design

Behavioral-Level Design: Start with C (or Similar Language)

- Replace first step of RTL design method by two steps
- Capture in C, then convert C to high-level state machine
- How convert from C to high-level state machine?

Step 1A: Capture in C
Step 1B: Convert to high-level state machine

|  | Step | Description |
| :---: | :---: | :---: |
|  | Capsure high-level | Describe the system's desired behavior as a high-level state max |
|  | state machine | The state machine consists of states and tramsitions. The state machine |
|  |  | are more than just Boolean operations on bit inputs and catputs. |
|  | Create a datapath | Create a datapath to carry out the data operations of the high-level state machine. |
|  | Connect she datapath to a controller | Connect the datapath to a controller block. Connect external Boolean inputs and outputs to the controller block. |
| $\frac{\overrightarrow{\hat{e}}}{\hat{5}}$ | Derive the controller's FSM | Convert the high-level state machine to a finite-state machine (FSM) for the controller, by replacing data operations with setting and reading of control signals to and from the datapath. |

RTL Design
Converting from C to High-Level State Machine

- Convert each C construct to
equivalent states and
transitions
- Assignment statement
$\square$ Becomes one state with assignment
- If-then statement
- Becomes state with condition check, transitioning to "then" statements if condition true otherwise to ending state
- "then" statements would also be converted to states



## RTL Design

Converting from C to High-Level State Machine

## - If-then-else

- Becomes state with condition check, transitioning to "then" statements if condition true, or to "else" statements if condition false


While loop statement

- Becomes state with condition check, transitioning to while loop's statements if true, then transitioning back to condition check


RTL Design
Simple Example of Converting from C to High-Level State Machine

\}-
(a)


o Simple example: Computing the maximum of two numbers

- Convert if-then-else statement to states (b)
- Then convert assignment statements to states (c)


## RTL Design

Example: Sum-of-Absolute-Differences $C$


