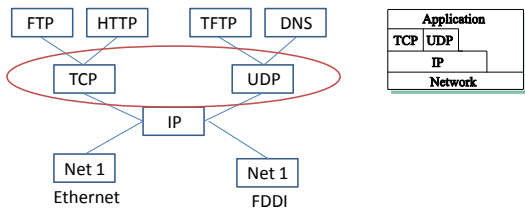


Fundamentals of Computer Networks ECE 478/578



Lecture #19: Transport Layer
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The Internet Architecture



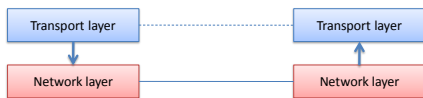
FTP: File Transfer Protocol TCP: Transmission Control Protocol
 HTTP: Hypertext Transport Protocol UDP: User Datagram Protocol
 TFTP: Trivial File Transfer Protocol IP: Internet Protocol
 DNS: Domain Name System

End-to-End Protocols

Responsible for the realization of host-to-host packet delivery

Services of the transport layer

- Guaranteed message delivery (end-to-end)
- Ordered delivery
- Rejection of duplicate messages
- Messages of arbitrary length
- Congestion control to handle network overloading
- Running of multiple application processes at the same host



Why do we need a Transport Layer?

Limitations of the network layer

- Can drop messages (routers may incur buffer overflow etc.)
- Can reorder messages
- Deliver duplicate copies of the same message
- Limit messages to finite size (eg. IP size is limited)

Network layer is intended to hide the different technologies on a link level and provide a transparent end to end routing service

Transport layer provides **end-to-end transport service** to the application layer

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Internet Transport Protocols

Datagram messaging service (UDP)

"Best-effort" IP end-to-end service

Reliable, in-order delivery (TCP)

- Connection set-up
- Discarding of corrupted packets
- Retransmission of lost packets
- Flow control
- Congestion control (What is the difference?)

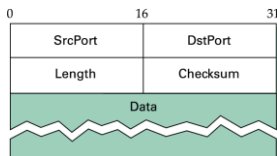
Remote Procedure Call (RPC)

Request/Reply Service

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User Datagram Protocol (UDP)

Process-to-process communication service

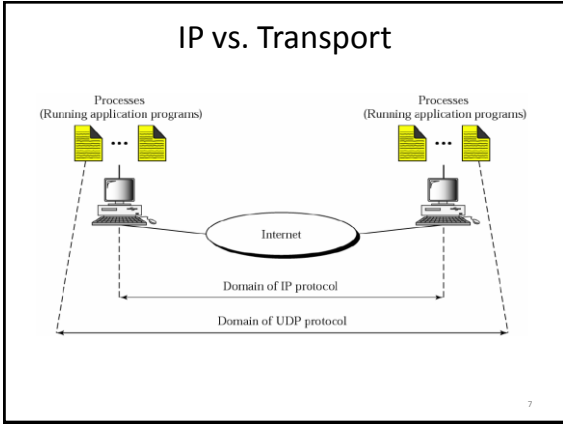


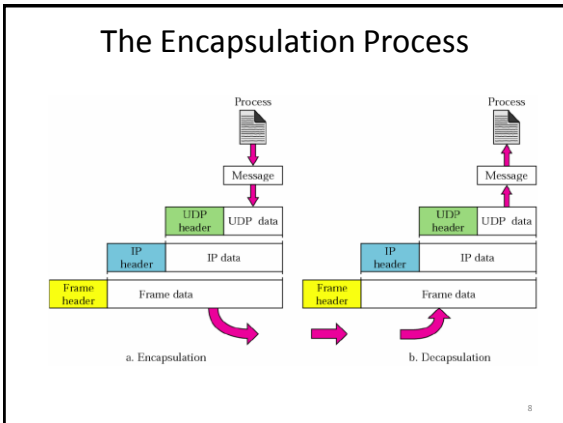
Processes are identified based on **incoming ports (sockets in Windows)**

16 bits for each field yields 64K different identifiers

<IP, port> combination allows de-multiplexing at receiving host

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Port Discovery

Use **well-publicized** ports for different services

- DNS uses port 53
- Email uses port 25
- HTTP uses port 80

Use one port as a "port-mapper" service

- Call 411 to learn the port of any other process
- Allows for dynamic allocation of ports to different services
- Allows for the assignment of ports to newly created services

Multiplexing and De-multiplexing

Host may be running multiple processes at the same time

These processes

- Generate multiple messages for the same host
- Generate multiple messages for multiple hosts

Transport layer multiplexing

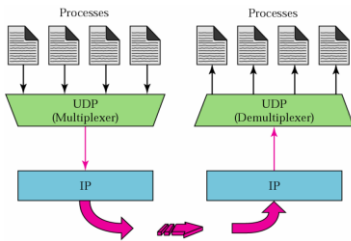
- Multiplex messages from multiple processes
- Break down messages to segments and pass to network layer

Transport layer de-multiplexing

- Reassemble messages at the receiving host and pass to the communication processes

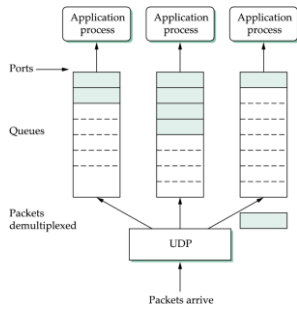
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Multiplexing and De-multiplexing

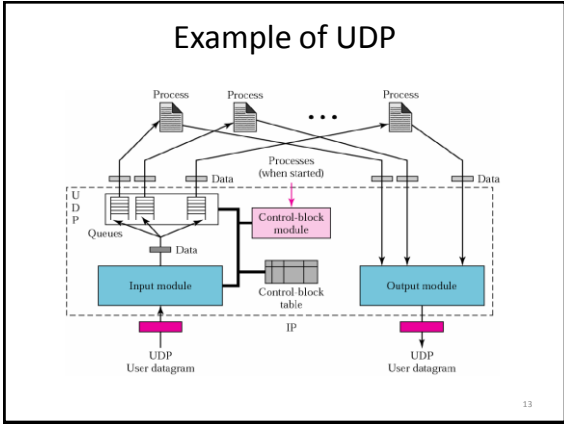


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UDP Message Queue



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Control Block Module

Responsible for the management of the Control Block Table
 Request for a port for a new starting process
 Update the control block table

Example

State	Process ID	Port Number	Queue Number
In Use	2,569	52,010	34
In Use	4,759	52,011	
Free	-	-	-
In Use	7,489	80	26
Free	-	-	-

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Input Module

Receive UDP packet from IP layer
 Look up control block table to map port
 If a queue exists, push packet to the corresponding queue
 If not, allocate queue for the new process

State	Process ID	Port Number	Queue Number
In Use	2,569	52,010	34
In Use	4,759	52,011	45
Free	-	-	-
In Use	7,489	80	26
Free	-	-	-

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Output Module

Receive data from the application layer
Create a UDP packet and send it.

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Advantages of UDP

Control over what data is sent and when

As soon as an application process writes into the socket
... UDP will package the data and send the packet

No delay for connection establishment

UDP just sends messages without contacting the host first
Pays off when host is expecting messages anyway

Stateless connection

No allocation of buffers, parameters, sequence #s, etc.
... making it easier to handle many active clients at once (think of servers)

Small packet header overhead

UDP header is only eight-bytes long

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Disadvantages of UDP

"Best effort" networking

No guarantee delivery of messages to destination host, no ordered delivery

No congestion control

No adaptation to the congestion conditions of the network

Suppresses TCP flows

In case of congestion TCP flows will back off while UDP will stay on the same rate
Can be used as an attack method (UDP flooding attack)

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Applications Utilizing UDP

Simple query protocols like Domain Name System

- Delay for connection establishment is too large
- Queries are small and UDP adds a small overhead (header)
- Easier to have application retransmit if needed
- Usually may fit within a UDP packet so no out-of-order danger

Multimedia Applications

- Retransmitting lost/corrupted packets is not worthwhile
- By the time the packet is retransmitted, it's too late
- E.g., telephone calls, video conferencing, gaming
- Certain loss is acceptable since Voice, picture, etc are still discernable

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